

Practical Session 2 : Foster Design

Objective

The objective of this practical session is to apply Foster design methodology to some specific problems.

1 – Maximum value

Suppose you are given a two-dimensional array of dimensions M and N containing real numbers. We make the assumption that M and N are much larger than the number of available processing units p . We also assume that the function $\max(a, b)$ is available.

Question 1

Discuss the first step in Foster methodology. Indicate how you make your choices. You may use diagrams to explain the reasoning.

2 – Matrix vector product

We consider the matrix-vector multiplication $y = Ax$, where $A \in \mathbb{R}^{n \times n}$, $x \in \mathbb{R}^n$ and $y \in \mathbb{R}^n$. This product is to be computed on a system with p processing units. In the first stage, we do not consider the topology of the network.

The matrix A has been partitioned column-block-wise in blocks A_i of size $n \times (n/p)$ and vector x in blocks X_i of size n/p respectively,

$$A = (A_1, A_2, \dots, A_p)$$

and

$$x = \begin{pmatrix} X_1 \\ X_2 \\ \vdots \\ X_p \end{pmatrix}$$

Question 2

Discuss the first step in Foster methodology. Indicate how you make your choices. You may use diagrams to explain the reasoning.

3 – Matrix product

We consider the following algorithm for the product of two matrices

Algorithm 1 : Matrices product

```

Data : Matrices A and B
Result : Matrix C
1 for  $i \leftarrow 1$  to  $n$  do
2   for  $j \leftarrow 1$  to  $p$  do
3     sum = 0
4     for  $k \leftarrow 1$  to  $m$  do
5       sum =  $A_{ik}B_{kj}$ 
6     end
7      $C_{ij} = \text{sum}$ 
8   end
9 end

```

Question 3

Design an algorithm that enable the computation of matrices product in parallel.

4 – Maximum subarray

In order to find the maximum subarray, that is find the subarray of a matrix with the maximum sum, we can use Kadane's algorithm algorithm

Algorithm 2 : 1D Kadane's algorithm

```

Data : array
Result : maxSum, left, right
1 (maxSum, left, right) =  $(-\infty, 0, 0)$ 
2 currentMaxSum = 0
3 currentStartIndex = 0
4 for  $i \leftarrow 1$  to  $n$  do
5   | currentMaxSum = currentMaxSum + array(i)
6   | if  $currentMaxSum > maxSum$  then
7   |   | (maxSum, left, right) = (currentMaxSum, currentStartIndex, i)
8   |   end
9   | if  $currentMaxSum < 0$  then
10  |   | currentMaxSum = 0
11  |   | currentStartIndex =  $i + 1$ 
12  |   end
13 end

```

The 2D algorithm relies on the use of Kadane's algorithm

Algorithm 3 : 2D maximum subarray

```

1 Data : array
Result : maxSum, left, right, top, bottom
2 (maxSum, left, right, top, bottom) =  $(-\infty, 0, 0, 0, 0)$ 
3 for  $i \leftarrow 0$  to  $n$  do
4   | temp(0 :n-1) = 0
5   | for  $j \leftarrow i$  to  $n$  do
6   |   | for  $k \leftarrow 0$  to  $m$  do
7   |   |   | temp(k) += array(j,k)
8   |   |   end
9   |   | sum = kadane(temp, start, finish)
10  |   | if  $sum > maxSum$  then
11  |   |   | (maxSum, left, right, top, bottom) = (sum, i, j, start, finish)
12  |   |   end
13  |   end
14 end

```

Question 4

Design an algorithm that enable the resolution of the maximum subarray problem in parallel in 2D.